

# Proving termination with transition invariants of height $\omega$

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# Termination Theorem by Podelski and Rybalchenko

- ▶ Transition invariants are used by Podelski and Rybalchenko to prove the termination of a program.
- ▶ A **transition invariant** of a program is a binary relation over program's states which contains the transitive closure of the transition relation of the program; i.e.  $T \supseteq R^+ \cap (\text{Acc} \times \text{Acc})$ .
- ▶ A relation is **disjunctively well-founded** if it is a finite union of well-founded relations.

## Theorem (Termination Theorem)

*The program  $P$  is terminating iff there exists a disjunctively well-founded transition invariant for  $P$ .*

# The proof by Podelski and Rybalchenko requires Ramsey Theorem

If you have  $\omega$  people at a party then either there exists an **infinite** subset whose members all know each other or an **infinite** subset none of whose members know each other.

## Theorem (Ramsey for pairs)

Assume  $I$  is a set having some injective enumeration

$\sigma = x_0, x_1, \dots, x_i, \dots$ . Assume  $S_1, \dots, S_n$  are binary relations on  $I$  which are a partition of  $\{(x_i, x_j) \in I \times I : j < i\}$ , that is:

1.  $S_1 \cup \dots \cup S_n = \{(x_i, x_j) \in I \times I : j < i\}$
2. for all  $1 \leq k < h \leq n$ :  $S_k \cap S_h = \emptyset$ .

Then for some  $k \in [1, n]$  there exists some infinite  $X \subseteq \mathbb{N}$  such that:

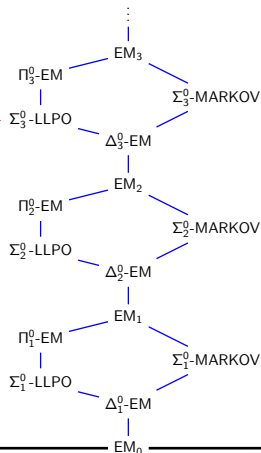
$$\forall i, j \in X. (j < i \implies x_i S_k x_j).$$

How many steps before the program  $P$  terminates? Hard to say, because Ramsey Theorem is a purely **classical** result.

Classical Logic

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$RT_2^2$



We provided a new intuitionistic version of Ramsey Theorem, which we called *H-closure Theorem*, and a new intuitionistic proof of the Termination Theorem.

## Why did we care?

In order to use our proof to extract some *effective bounds* for the Termination Theorem.

# H-closure Theorem

In order to analyse the bounds implicit in the Termination Theorem we replaced Ramsey with some intuitionistic result.

Let  $R$  be a binary relation on  $I$ .

- ▶  $H(R)$  is the set of the  **$R$ -decreasing transitive finite sequences** on  $I$ :

$$\langle x_1, \dots, x_n \rangle \in H(R) \iff \forall i, j \in [1, n] (i < j \implies x_j R x_i).$$

- ▶  $R$  is  $H$ -well-founded if  $H(R)$  is  $\succ$ -well-founded.

## Theorem (H-closure)

*The H-well-founded relations are closed under finite unions:*

$$(R_1, \dots, R_n \text{ H-well-founded}) \implies ((R_1 \cup \dots \cup R_n) \text{ H-well-founded}).$$

# Our goal

## Theorem

*A function has at least one implementation which has a disjointively well-founded transition invariant whose relations are primitive recursive and have height  $\omega$  if and only if it is primitive recursive.*

# Our goal

## Theorem

Assume that  $P$  is such that  $R^+ \cap (\text{Acc} \times \text{Acc}) = R_1 \cup \dots \cup R_k$ , where

1. the complement of the set of the final states of  $R$  ( $S \setminus F$ ) is a primitive recursive set;
2.  $R$  is the graph of a primitive recursive function  $t : S \setminus F \rightarrow S$ : i.e.

$$R = \{(x, t(x)) : x \in S \setminus F\}.$$

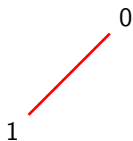
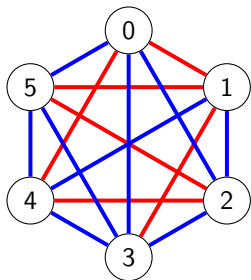
3.  $R_1, \dots, R_k$  are primitive recursive relations and have height  $\omega$ .

Then there exists a primitive recursive  $g : S \rightarrow \mathbb{N}$  such that  $t^{g(s)}(s) = t^{g(s)+1}(s)$ .



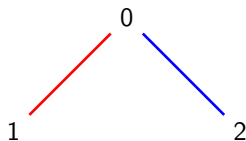
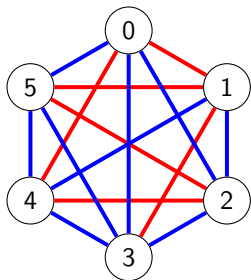
# Erdős' trees

Assume given a sequence  $\langle 0, \dots, 5 \rangle$  such that the coloring between its elements is as follows.



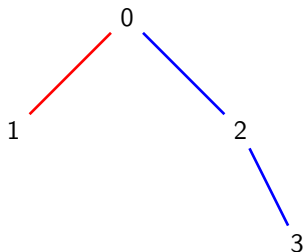
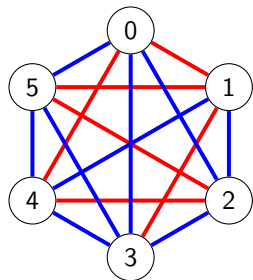
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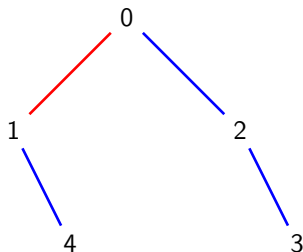
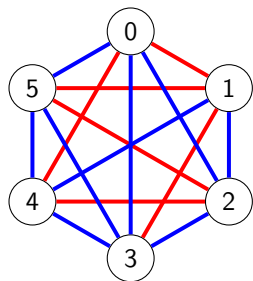
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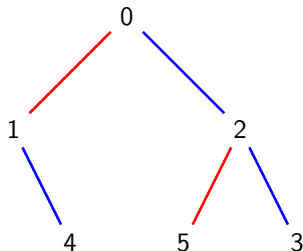
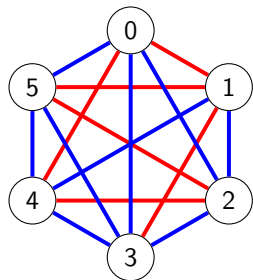
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## Sketch of the proof

- ▶ There exists a **primitive recursive function** from the set of  $k$ -branching Erdős' trees to  $\omega^k$ , which we called  $h_k$ , which “preserves” the **one step expansion** between trees:

$$T' \succ_E T \implies h_k(T') < h_k(T).$$

- ▶ There exists a **primitive recursive function**  $E$  from  $H(R_0 \cup \dots \cup R_{k-1})$  to the set of the Erdős' trees (as we did in the previous example), which “preserves” the **one step expansion** between sequences:

$$L' \succ L \implies E(L') \succ_E E(L).$$

So given an  $R$ -chain  $\alpha$  we want to prove that it is finite.

- ▶ Define

$$\begin{aligned}\theta : \mathbb{N} &\rightarrow H(R_0 \cup \dots \cup R_{k-1}) \\ m &\mapsto \langle s_0, t(s_0), \dots, t^m(s_0) \rangle\end{aligned}$$

By definition we have:

$$m' > m \implies \theta(m') \succ \theta(m).$$

- ▶ Let  $\phi : \mathbb{N} \rightarrow \mathbb{N}^k$  be such that  $\phi(m) = h_k(E(\theta))$ . It is **primitive recursive!**

## Lemma

For each  $\sigma : \mathbb{N} \rightarrow \mathbb{N}^k$  primitive recursive, there exists  $g : \mathbb{N} \rightarrow \mathbb{N}$  primitive recursive such that

$$\forall n \exists m \in [n, g(n)] (\sigma(m) \preceq_k \sigma(m+1)).$$

By applying this lemma to  $\phi$  we obtain the primitive recursive **bound** we were looking for!









# What are we doing now?

Since by Termination Theorem Podelski, Rybalchenko and Cook produced an algorithm, we hope that  $H$ -closure Theorem could be useful in order to find **bounds** also for the algorithm.

## Conjecture

*A function has at least one implementation in Podelski-Rybalchenko language which the Terminator Algorithm may catch terminating if and only if the function is primitive recursive.*

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